# Modal logics of polytopes – what we know so far

David Gabelaia

in collaboration with Members of Esakia Seminar

Guram Bezhanishvili, Nick Bezhanishvili, Mamuka Jibladze, Evgeny Kuznetsov, Kristina Gogoladze, Maarten Marx, Levan Uridia et alii

# Topology and modal logic

#### McKinsey and Tarski 1944

- Interpret propositions as subsets of a topological space
- Interpret Boolean operations as their set-theoretic counterparts
- Interpret the modal diamond as closure, or as derivative
- S4 is the modal logic of any crowded, separable, metrizable space
- Rasiowa and Sikorski 1963
  - S4 is the modal logic of any crowded, metrizable space
- $\blacksquare$  So any  $\mathbb{R}^n$  generates S4









# Map of an Island Mapping f S Α В







# (M)Any subsets – wild logics

• Any finite connected quasiorder (S4-frame) is an interior image of  $\mathbb{R}^n$ 

[G. Bezhanishvili and Gehrke, 2002]

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# (M)Any subsets – wild logics

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- The subalgebras of the closure algebra ( (R<sup>n</sup>), C) generate all connected extensions of S4
- The subalgebras of the closure algebra ( (Q), C) generate all normal extensions of S4

[G. Bezhanishvili, DG and Lucero-Bryan, 2015]

• Too many subsets!

## Nice subsets – tame logics?

• Piecewise linear subsets = polytopes



#### Catmull-Clark









## Nice subsets – tame logics?

• Piecewise linear subsets = polytopes

**PC**<sup>n</sup> = C-logic of all polytopal subsets of  $\mathbb{R}^n$ **PD**<sup>n</sup> = d-logic of all polytopal subsets of  $\mathbb{R}^n$ 

Our aim is to investigate these modal systems

• In this talk - PC<sup>n</sup>

### General observations

If  $A \cap B = \emptyset$  and  $A \subseteq CB$ Then dim(A) < dim(B)

Put  $\beta A \equiv CA \setminus A$  (boundary of A) Then  $\beta^n A = \emptyset$  iff dim(A) < n

It follows that each **PC**<sup>n</sup> is a logic of finite height.

## Forbidden frames for **PC**<sup>n</sup>



## Forbidden frames for **PC**<sup>n</sup>



## Forbidden frames for **PC**<sup>n</sup>



• **PC<sup>1</sup>** is the modal logic of a 2-fork

[van Benthem, G. Bezhanishvili and Gehrke, 2003]



## PC<sup>2</sup> – forbidden frames



## **PC**<sup>2</sup> – forbidden frames



## PC<sup>2</sup> – forbidden frames



Any other forbidden configurations?

# Example



# Example



## Example





# PC<sup>2</sup> – admitted frames

<u>Lemma</u>: Any crown frame is a partial polygonal interior image of the plane.



## PC<sup>2</sup> – Axiomatization

#### Bad, but almost good guys

#### Very nice guys





# PC<sup>2</sup> – admitted frames

<u>Lemma:</u> Any rooted poset not reducible to any of the forbidden frames is a p-morphic image of a crown frame.

<u>Theorem</u>: The logic **PC**<sup>2</sup> is axiomatizable by Jankov-Fine axioms of the five forbidden frames.

## **PC<sup>3</sup>** – forbidden frames



## PC<sup>3</sup> – forbidden frames



#### Any other forbidden configurations?

# PC<sup>3</sup> – Spherical (open) polyhedra



# Planar graphs

 A graph is planar if it can be drawn on the plane (=on a surface of a sphere) without intersecting edges



## Non-planar graphs





**K**<sub>5</sub>

K<sub>3,3</sub>

## **PC<sup>3</sup>** – forbidden frames



## PC<sup>3</sup> – forbidden frames



#### Anything else?

#### Face posets of sphere triangulations



